## Scalable and Secure Electronic Health Record (EHR) System in BigData Environment

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#### Abstract

Cloud computing services offer several flavors of Virtual Machines (VMs). As a result, the centralized approaches are difficult in handling large-scale data sets well on cloud using just one single VM even if the VM has the highest computation and storage capability. Maintaining the linkage structure and updating the statistic information will be huge when the data sets become large. Centralized approach probably suffers from low efficiency and scalability when handling large data sets. All data processed should fit in memory for centralized approaches, but theassumption fails to hold data-intensive on cloud. A large memory should be implemented in order to hold the records and registrations, but the total process is not fit in the centralized system. Also the scalability of the system is not up to the mark. To implement all the features there should be large memory. The centralized Top Down Specialization (TDS) approaches exploits the data structure Taxonomy Indexed Partition Structure (TIPS) to improve the scalability and efficiency by indexing anonymous data records and retaining statistical information in TIPS. The data structure speeds up the specialization process because with the help of indexing structure that helpsin avoiding frequently scanning entire data sets and storing statistical results, thus, thwart recomputation overheads. The amount of metadata that is retained for maintain the statistical information and other linkage related information of a record partitions is comparatively huge set. This consumes considerable amount of memory as like the original data set. The overhead incurredby maintain and updating these linkage structure, and the statistical information maintained will also be huge and becomes a larger data set. Hence, centralized approaches probably suffer from low scalability and efficiency while handling large-scale data sets.

Keywords: big data, scalable, TIPS, TDS, electronic health record

#### **1. INTRODUCTION**

Cloud computing is one of the most pre-dominant paradigm in latest traits for computing and storing purposes. It poses a significant impact on the current IT industry as well as on research communities

[1] [2] [3]. Data safety and privateness of information is one of the predominant situation within the cloud computing. Data anonymization has been notably studied and broadly followed approach for privacy preserving in information publishing and sharing methods. Data anonymization is displaying up of sensitive information for owner's information document to mitigate unidentified risk. The privacy data of person may be correctly maintained even as a few mixture records are shared to information person for information evaluation and information mining.

The rest of the paper is organized as follows. Section 2 discusses the related works, section 3 brings out the proposed work, section 4 gives the results and section 5 conclusion.

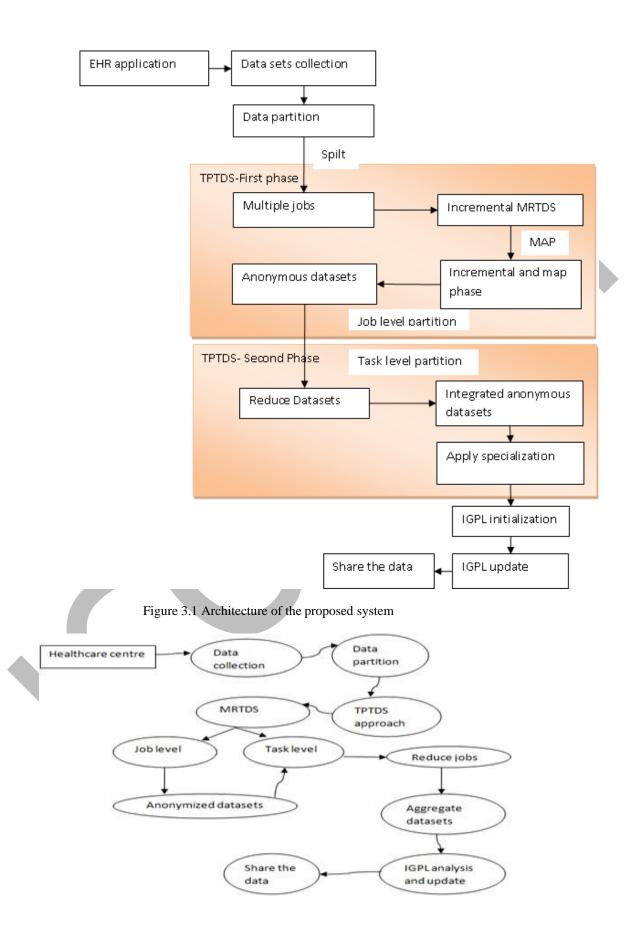
#### 2. RELATED WORKS

One of the extensively investigated concept in recent days is data privacy preservation [4]. The scalability problems of anonymization are addressed by introducing sampling techniques, decision trees, etc. Iwuchukwu and Naughton [5] proposed R-tree index-based approach by creating a spatial index over data sets by achieving high efficiency. Fung et al. [6], [7], [8] proposed a TDS approach, which aimed at producing anonymous data sets by not taking the data exploration problem [6]. In literature, we could find various distributed algorithms, which are proposed for preserving privacy of multiple data sets retained by multiple parties. Mohammed et al. [9] and Jiang [10] proposed distributed algorithms for anonymizing vertically partitioned data from different sources. This work aimed at preserving the privacy information without disclosing privacy information from one party to another. Mohammed et al. [11] and Jurczyk [12] proposed a distributed algorithmic approach to anonymize horizontally partitioned data set.

#### **3. PROPOSED SYSTEM**

Figure 3.1 shows the architecture diagram of the proposed system. The health records of various patients that have been registered are stored in the database. Generally, big data can be processed incrementally and data are divided into equally sized blocks (data chunks). Data chunks created using k-anonymity and generate various anonymization levels. Figure 3.2 shows the data flow diagram of the proposed system. The data is collected from the healthcares and it is partitioned into data blocks. The partitioned data enter into the TPTDS phase with the help of Map-Reduce Top down Specialization (MRTDS) approach. The results of the two phases are the anonymized datasets. Then the datasets enters to the reduce phase to providecounts to the datasets. After that, the datasets are sent to the IGPL analysis and IGPL update phase to calculate the anonymization level. At last the data is shared

and accessed in the cloud.



#### 3.1. Module Description

There are four modules in the proposed system and they are as follows:

- a. EHR system development
- b. Data chunk creation
- c. MRTDS implementation
- d. IGPL analysis
- a) EHR System Development

EHRs are real-time, patient-centric management system that helps in maintaining records of patients that make information available instantly and securely to authorized users. Figure 3.3 shows the homepage design of EHR system. It contains various tabs named home, hospital, patient, research centre, laboratory, specialist, pharmacy. By clicking on those tabs, the details of the particular option can be obtained as well as entered. The design of the homepage contains the Administrator, who has the central control access of all the information in the system.

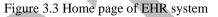
EHRs may include a range of data, including medical history, medication and allergies, immunization status, laboratory test results, radiology images, vital signs, personal statistics like age and weight, and billing information. EHR system is created to gather big data from the users of the system such as patients, doctors and health care's. The system is designed in such a way for representing data and that accurately captures the state of the patient at all times. It allows for an entire patient history to be viewed without the need to track down the patient's previous medical record volume and assists in ensuring data is accurate, appropriate and legible. It cut down the timeinvolved in data replication as there is only one modifiable file, which means the file is constantly up to date, when viewed at a later date reduces the issue of paperwork. Due to all the information being maintained in a single file, it makes it much more effective when extracting medical data for the examination of possible trends and long term changes in the patient condition.

While an EHR does contain the medical and treatment histories of patients, it is built to go beyond standard clinical data collected in a provider's office and can be inclusive of a broaderview of a patient's care. One of the prime features of an EHR is that, the authorized providers can create and manage the information in a digital format that is capable of being shared with other providers across more than one health care organization. It can share information with other health care providers and organizations – such as laboratories, specialists, medical imaging facilities, pharmacies, emergency

The algorithm for data partition map and reduce is as follows:

facilities, and school and workplace clinics - so they contain information from all clinicians involved in a patient's care.

WELCOME TO HEALTHCARE SYSTEM			
HOME	HOSPITAL	PATIENT	RESEARCH CENTER
ADMINISTRATOR USER NAME PASSWORD			



#### b) Data Chunk Creation

The data records available in the file or database are spitted into data chunks. A chunk is a fragment of information. Each chunk contains a header which indicates some parameters e.g. the type of chunk, size. In the middle, there is a variable area, containing data which are decoded by the program from the parameters in the header. Chunks may also be fragments of information which are downloaded or managed by distributed programs. In distributed computing, a chunk is a set of data which are sent to a processor or one of the parts of a computer for processing. Figure 3.4 shows the creation of data chunks.

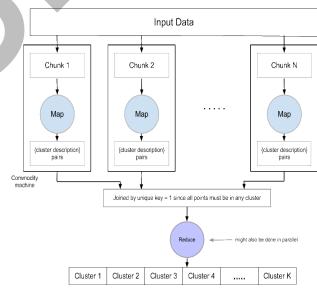


Figure 3.4 Data chunk creation

list(r)).

Once partitioned data sets  $D_i$ ,  $1 \le i \le p$  are obtained, we run MRTDS ( $D_i$ ,KI,AL<sup>0</sup>) on these datasets in parallel

to derive intermediate levels  $AL_i^*$ ,  $1 \le i \le p$ .

AL- Anonymization level

To employ anonymisation level and to manage the process of anonymisation, the followingsteps are to be followed:

Step 1 : Initializes the values of information gain and privacy loss for all specializations,

with the job IGPL initialization.

Step 2 : i) It is an iterative process, which is the main body of the algorithm.

ii) The best specialization is selected from valid specializations in current anonymisationlevel as described in step i.

iii) A specialization is said to be valid, if it satisfies two conditions. One is that its parent value is not a leaf, and the other is that the anonymity  $A_c(spec) > k$ , i.e., the data set is still k- anonymous if is performed. Then, the current anonymisation level is modified via performing the best specialization in Step ii, i.e., removing the old specialization and inserting new ones that are derived from the old one. In Step iii, information gain of the newly added specializations and privacy loss of all specializations need to be recomputed, which are accomplished by the job IGPL Update. The iteration continues until all specializations become invalid. Ultimately, D is concretely anonymised according to the resultant AL in Step 3 by the job Data Anonymisation.

MRTDS produces the same anonymous data as the serial TDS, because they follow the same steps. MTRDS mainly differs from the serial TDS in calculating IGPL values. However, calculating IGPL values dominates the scalability of TDS, as it requires TDS algorithms to count the statistical information of data sets iteratively. MRTDS exploits Map Reduce to carry out the computation of IGPL in a parallel and scalable way.

- i- Denotes the count of the data to be portioned
- ii- Di Partitioned Data

D - Data set

c) MRTDS Implementation

MRTDS implements map reduce system which contains two steps, such as map andreduce. Incremental map provides the necessary control over the alignment and granularity of the input parts. Incremental reduce processes the output of the map function grouped by the keys of the generated key value pairs. Map reduce version of MRTDS concretely conducts the computation required in TPTDS. MRTDS only leverages the task level parallelization's of map reduce. The map function emits anonymous records. The reduce function simply aggregates these anonymous records and counts their number. The figure 3.5 shows that the execution framework overview of MRTDS.

The algorithm for MRTDS driver is as follows:

Input : Data set D, the topmost AL  $AL^0$ , and k-anonymity parameter k. Output : Anonymous data set  $D^{*}$ .

Step 1 : Initialize the IGPL values for each specialization spec  $\Box U^m j=1$  .Cutj with respect toAL<sup>0</sup>, using the job IGPL initialization;

Step 2 : While spec  $\Box$  U <sup>m</sup>j=1 Cut j is valid.Find the best specialization from AL<sub>i</sub>, specBest;

#### i. Update $AL_i$ to $AL_{i+1}$ ;

ii. Update information gain of the new specializations in AL<sub>i+1</sub>, and privacy loss foreach specialization via job IGPL Update;

Step 3: Anonymise data set D in terms of resultant AL via job DataAnonymisation.

d) IGPL Analysis

The main task of IGPL Initialization is to initialize information gain and privacy loss of all specializations in the initial AL. The statistical information is required for each specialization to calculate information gain. The number of records in each current QI-group needs to be computed. The algorithm for IGPL Initialization Map is as follows:

Input : Data record (ID<sub>r</sub>,r),  $r\Box D$ ; anonymisation level AL. Output: Intermediate key-value pair (Key,Count)

Step 1: For each attribute value v<sub>j</sub> in r, find its

specialization in current AL:spec; Let q be theparent in spec, and c be the q's child value that is also ancestor of v<sub>i</sub> in TT<sub>i</sub>;

Step 2: For each vi, emit ((q, c, sv), count);

Step 3: Construct quasi-identifier

qid\*=(q1,q2...qm) where qi,  $1 \le i \le$ m, is the parent of specialization in current AL. Emit

((qid\*,\$,#),count);

Step 4: For each  $i\square[1,m]$ , replace qi in qid\* with its child ci, where ci is also the ancestor of vi.

Let the resultant quasi identifier be qid. Emit ((qid\*,pi,#),count).

The inputs are data sets that consist of a number of records. It is the sequence number of the record. Step 1 gets the potential specialization for the attribute values. Then Step 2 emits key- value pairs containing the information of specialization, sensitive value, and the count information of this record. According to the above information, we compute information gain for a potential specialization in the corresponding Reduce function. Step 3 aims at computing the current anonymity, while Step 4 is to compute anonymity after potential specializations. The symbol '#' is used to identify whether a key is emitted to compute information gain or anonymity loss, while the symbol '\$' is employed to differentiate the cases whether a key is for computing Aa(spec) or Ac(spec).

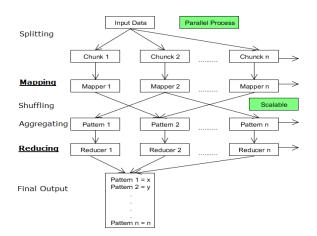


Figure 3.5 Execution framework overview of MRTDS

The algorithm for IGPL Initialization Reduce is as follows: Input : Intermediate key-value pair (key, list (count)).

Output: Information gain

(spec,IG(spec) ) and anonymity (spec,Ac(spec))), (spec,Aq(spec))) for all specializations

Step 1 : For each key, sum count;

Step 2 : For each key, if key.sv≠#, update statistical counts:

- i.  $|(R_c,sv)| \square sum$ ,  $|R_c| \square sum + |R_c|, |(R_q,sv)| \square sum + |(R_q,sv)|, |R_q|$  $\square sum + |R_q|$
- ii. If all sensitive values for child have arrived, compute  $I(R_c)$  according to (4-3);
- iii. If all children c for parent p have arrived, compute I(Rq)and IG(spec), Emit (spec

,IG(spec));

Step 3: For each key, if key.sv=#, update anonymity.

- i. If key.c=\$ and sum<Aq(spec), update current anonymity:Aq(spec) sum;
- ii. If key.c≠\$ and sum<A<sub>c</sub>(spec), update potential anonymity ofSpec: A<sub>c</sub>(spec) □sum;

Step 4: Emit (spec, Aq(spec)) and emit (spec, Ac(spec)).

The first step is to accumulate the values for each input key. If a key is for computing information gain, then the corresponding statistical information is updated in Step 2(i) I(Rq), I(Rc) and IG(spec) are calculated if all the count information they need has been computed in Step 2 (ii and iii) in terms of (3) and (4). A salient Map Reduce feature that intermediate key-value pairs are sorted in the shuffle phase, makes the computation of IG(spec) sequential with respect to the order of the specializations arriving at the same reducer. Hence, the reducer just needs to keep statistical information for one specialization at a time, which makes the reduce algorithm highly scalable.

To compute the anonymity of data sets before and after a specialization, Step 3 (i) finds the smallest number of records out of all current QI-groups, and Step 3 (ii) finds all the smallest number of records out of all potential QIgroups for each specialization. Step 4 emits the results of anonymity. Note that there may be more than one keyvalue pair (spec, A(spec)) for one specialization in output files if more than one reducer is set. But we can find the smallestanonymity value in the driver program. Then the privacy loss PL(spec) is computed. Finally, IGPL(spec) for each specialization is obtained.

The algorithm for IGPL update map as follows:

Input : Data record (IDr, r),  $r \Box D$ ; Anonymisation level AL.

Output: Intermediate key-value pair (key,count).

Step 1: Let attr be the attribute of the last best specialization. The value of this attribute in r is v. Find its specialization in current AL: spec. Let q be the parent in spec, and c be q's child that is also an ancestor of v; Emit((q,c,sv),count); Step 2: Construct quasi-identifier  $qid^* = (q1, q2, q2)$ .....

qm), qi ,1 $\leq$ i $\leq$ m, is the

parent of a specialization in current AL and is also an ancestor of viin r. Step 3: For each i □[1,m], replace qi in qid<sup>\*</sup> with its child cj if the specialization related to qi is valid, where cj is also the ancestor of vj.Let the

#### resultant quasi-



identifier be qid. Emit ((qid, qi, #),count).

The IGPL Update job dominates the scalability and efficiency of MRTDS, since it is executed iteratively as described in MRTDS driver. The IGPL Update job is quite similar to IGPL Initialization, except that it requires less computation and consumes less network bandwidth. Thus, the former is more efficient than the latter. IGPL update map describes the Map function of IGPL Update. The Reduce function is the same as IGPL Initialization, already described in IGPLinitialization reduce.

After a specialization spec is selected as the best candidate, it is required to compute the information gain for the new specializations derived from spec. So, Step 1 in IGPL update map only emits the key-value pairs for the new specializations, rather than all in IGPL initialization map. Note that it is unnecessary to re-compute the information gain of other specializations because conducting the selected specialization never affects the information gain of others. Compared with IGPL initialization, only part of data is handled and less network bandwidth is consumed.

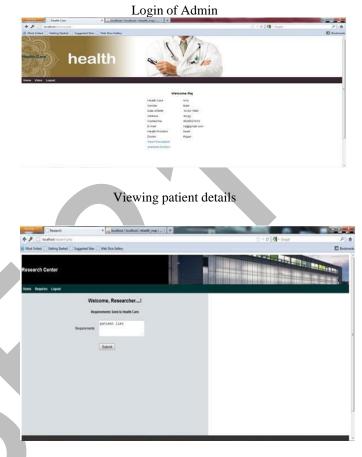
On the contrary, the anonymity values of other specializations will be influenced with high probability because splitting QI-groups according to spec changes the minimality of the smallest QI-group in the last round. Therefore, it is needed to compute  $A_c(spec)$  for all specializations in AL, as described in Step 2 and 3 of IGPL update map. Yet Aq (spec) can be directly obtained from the statistical information kept by the last best specialization. Note that if the specialization related to qi in Step 3 is not valid, no resultant quasi-identifier will be created.

Since the IGPL update job dominates the scalability and efficiency of MRTDS, hence analyze its complexity as follows. Let n denote all the records in a data set, m be the number of attributes, s be the number of mappers and t be the number of reducers. As a mapper emits (m+1) key-value pairs, it takes O(1) space and O(m\*n/s) time a reducer takes O(1) space and O(m\*n/s)

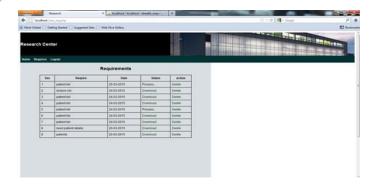
time. Note that a reducer only needs O(1) space due to the Map Reduce feature that the key value pairs are sorted in the shuffle phase. Otherwise, the reducer needs more space to accumulate statistic information for a variety of specializations. The communication cost is O(m\*n) according to the map function, but communication traffics

can be reduced heavily by Map Reduce techniqueslike Combiner.

### 4. **RESULTS**



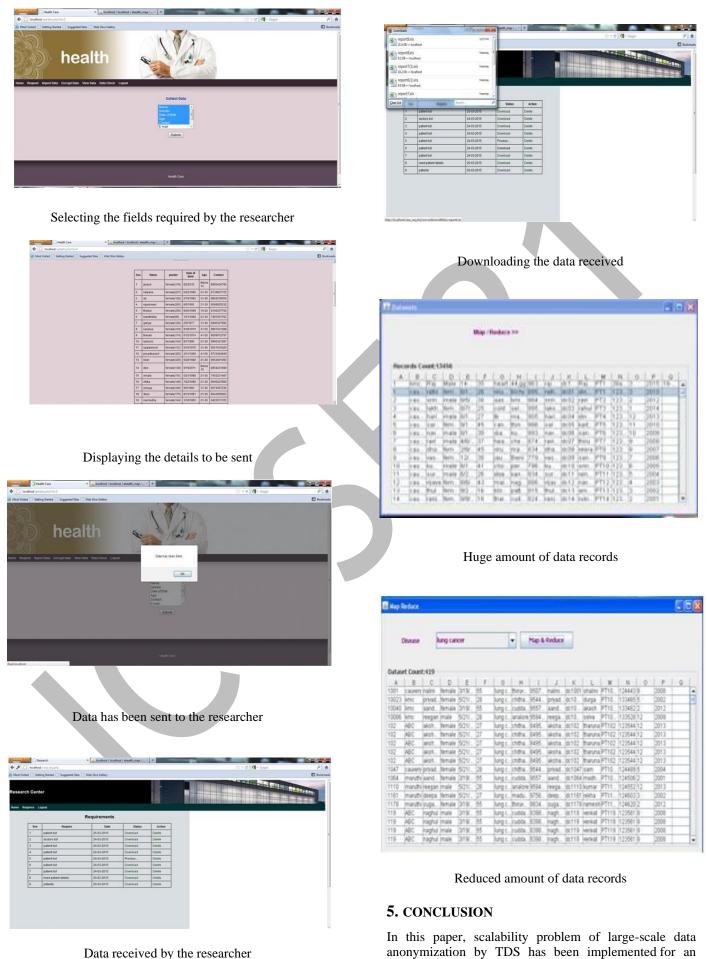
Raising query to the admin by researcher



Processing of the query

Responding to the query by the admin





Data received by the researcher

electronic health record maintenance application. A highly scalable two-phase TDS approach using MapReduce on cloud has been proposed. In the first phase, data sets were partitioned and anonymized. In the second phase, the intermediate results, generated by the first phase were merged and further anonymized for producing consistent k-anonymous data sets. In cloud environment, the privacy preservation for data analysis, share and mining is a challenging research issue due to increasingly larger volumes of data sets, thereby requiring intensive investigation. Based on the contributions herein, we plan to further explore the next step onscalable privacy preservation aware analysis and scheduling on large-scale data set.

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